

Course Name: Production & Operations Mgmt.

Unit: II → Operation Research-II

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Transportation Model Problems: →

Introduction: In the previous semester, the general nature of LPP & its solution (graphical, simplex, etc.) was discussed. However, as the number of variables & constraints increase, the computation becomes more difficult. One such model requiring simplified calculations is the distribution model or the transportation model.

Transportation problem is one of the sub-classes of LPP in which the main objective is to transport various quantities of or single homogenous commodity that are initially stored at various origins, to different destinations in such a way that total transportation cost is minimized. To achieve this objective, we must know the amount & location of available supplies, the quantity demanded, in addition to the involvement of cost associated with each sending.

Definition of T. Model:

- Feasible Solution: A feasible soln. to a transportation problem is a set of non-negative allocations (units transported per route from origin to destination),  $x_{ij}$  that satisfies the  $x_{ij}$  (row & column) restrictions.
- Basic Feasible Solution: A soln that contains not more than  $m+n-1$  non-negative allocations, where  $m$  is the no. of rows &  $n$  is the no. of columns of T. problem.
- Optimal Solution: A feasible soln (not necessarily basic) that minimizes (maximizes) the transportation cost (profit) is called an Optimal solution.



# Tabular Representation of T. Model :-

ORIGIN / SOURCE <small>Site</small>	DESTINATIONS					Supply / Availability
$i \downarrow j \rightarrow$	$d_1$	$d_2$	$d_3$	$d_4 \dots$	$d_n$	
$O_1$	$C_{11}$ $x_{11}$	$C_{12}$ $x_{12}$	$C_{13}$ $x_{13}$	$C_{14}$ $x_{14}$	$C_{1n}$ $x_{1n}$	$S_1$
$O_2$	$C_{21}$ $x_{21}$	$C_{22}$ $x_{22}$	$C_{23}$ $x_{23}$	$C_{24}$ $x_{24}$	$C_{2n}$ $x_{2n}$	$S_2$
$\vdots$						
$O_m$	$C_{m1}$ $x_{m1}$	$C_{m2}$ $x_{m2}$	$C_{m3}$ $x_{m3}$	$C_{m4}$ $x_{m4}$	$C_{mn}$ $x_{mn}$	$S_m$
Requirement / Demand	$b_1$	$b_2$	$b_3$	$b_4$	$b_n$	$\sum_{i=1}^m S_i = \sum_{j=1}^n b_j$

$C_{11}$  is the transportation cost from origin 1 to destination 1.  
 $x_{11}$  is the allocation (no. of units transported per route from origin 1 to dest. 1)  
 $b_1, b_2 \dots b_n$  is the requirement at respective destinations.

## Methods used to Solve Transportation Problems:

- 1) North-West Corner Method
- 2) Row Minima Method
- 3) Column Minima Method
- 4) Lowest Cost or Matrix Minima Method.
- 5) Vogel's Approximation Method (VAM) or Penalty Method.
- 6) Modified Distribution Method or MODI Method or U-V Method.

Solution rules to above methods on book (OR by AP Verma).

N<sub>11</sub> Solve the following problem using NWCM, RMM, CMM & MMM. (above 4 methods).

Origin	$d_1$	$d_2$	$d_3$	$d_4$	Supply
$O_1$	21	16	25	13	11
$O_2$	17	18	14	23	13
$O_3$	32	27	18	41	19
Demand	6	10	12	15	

$\sum \text{Supply} = \sum \text{Demand}$   
 $43 = 43$



## NORTH - WEST CORNER METHOD:

Origin	$d_1$	$d_2$	$d_3$	$d_4$	Supply
$O_1$	21 <sup>16</sup>	16	25	13	115
$O_2$	17	18	14	23	13
$O_3$	32	27	18	41	19
Demand	6	10	12	15	

Origin	$d_2$	$d_3$	$d_4$	Supply
$O_1$	16 <sup>5</sup>	25	13	5
$O_2$	18	14	23	13
$O_3$	27	18	41	19
Demand	5	12	15	



Origin	$d_2$	$d_3$	$d_4$	Supply
$O_2$	18 <sup>5</sup>	14 <sup>8</sup>	23	13
$O_3$	27	18 <sup>4</sup>	41 <sup>15</sup>	19
Demand	5	4	15	

$m = \text{no. of rows} = 4$   
 $n = \text{no. of columns} = 3$   
 $m+n-1 = 6$  reqd. no. of allocations.

$\therefore$  Actual no. of allocations after solving = 6.  
 $\therefore$  we have got basic feasible soln,  $\therefore$  is optimal.

Total cost =  $\sum (\text{allocation} \times \text{total cost})$   
 $= C_{11}x_{11} + C_{12}x_{12} + C_{22}x_{22} + C_{23}x_{23} + C_{33}x_{33} + C_{34}x_{34}$

$= 21 \times 6 + 16 \times 5 + 18 \times 5 + 14 \times 8 + 18 \times 4 + 41 \times 15$   
 $= \underline{\underline{\text{Rs } 1095}}$

## → Row - MINIMA METHOD:

Origin	$d_1$	$d_2$	$d_3$	$d_4$	Supply
$O_1$	21	16	25	13 <sup>11</sup>	11
$O_2$	17	18	14 <sup>12</sup>	23	13
$O_3$	32	27 <sup>10</sup>	18	41 <sup>4</sup>	19
Demand	6	5	10	12	

No. of allocations  
 $= 6$

$\therefore$  Total Cost =  $13 \times 11 + 17 \times 1 + 14 \times 12 + 32 \times 5 + 27 \times 10 + 41 \times 4$   
 $= \text{Rs } 922$



### 3) Column Minima Method :

Origin	$d_1$	$d_2$	$d_3$	$d_4$	Supply
$D_1$	21	16 <sup>10</sup>	25	13 <sup>1</sup>	11
$D_2$	17 <sup>6</sup>	18	14 <sup>7</sup>	23	13 <sup>7</sup>
$D_3$	32	27	18 <sup>5</sup>	41 <sup>14</sup>	19 <sup>14</sup>
Demand	6	10	12 <sup>5</sup>	15 <sup>14</sup>	

$$\text{Total Cost: } 17 \times 6 + 16 \times 10 + 14 \times 7 + 18 \times 5 + 13 \times 1 + 41 \times 14 \\ = \text{Rs } 1037$$

### 4) Matrix Minima Method

Origin	$d_1$	$d_2$	$d_3$	$d_4$	Supply
$D_1$	21	16	25	13 <sup>11</sup>	11
$D_2$	17 <sup>1</sup>	18	14 <sup>12</sup>	23	13 <sup>1</sup>
$D_3$	32 <sup>5</sup>	27 <sup>10</sup>	18	41 <sup>14</sup>	19 <sup>14</sup>
Demand	5	10	12	15 <sup>4</sup>	

$$\text{Total Cost: } 13 \times 11 + 17 \times 1 + 14 \times 12 + 32 \times 5 + 27 \times 10 + 41 \times 4 \\ = \text{Rs } 922$$

### 5) VAM (Imp.)

Here, we have to find penalties. Penalty means difference between lowest cost & next lowest cost. After determining penalty, we select the max. penalty in row & column. Then the lowest cost cell is taken in subsequent row or column & allocated.

Case: When there is tie in penalty between available & reqd., then penalty is shown having max. allocation. It is preferred to choose the cost difference corresponding to  $\infty$  largest no. of units can be assigned or corresponding to  $\infty$  cell chosen has minimum cost.



5

	$d_1$	$d_2$	$d_3$	$d_4$	Available	Row Penalties
$O_1$	1	2	1	4	30	$1-1=0$ ,
$O_2$	3	3	2	1	<del>50</del> 40	$2-1=0$ ,
$O_3$	4	2	5	9	20	$4-2=0$ ,
Reqd.	20	40	30	10		
Column Penalties	$3-1=2$	$2-2=0$	$2-1=0$	$4-1=3$		

	$d_1$	$d_2$	$d_3$	Available	R.P
$O_1$	1	<del>2</del> 10		<del>30</del> 10	$1-1=0$ , $O_1$ , 1
$O_2$	3	<del>3</del> 10	2	<del>40</del> 10	$3-2=1$ , 1 , 1
$O_3$	4	<del>2</del> 20	5	<del>20</del>	$4-2=2$
Reqd.	20	<del>40</del> 20	30		
C.P	$3-1=2$ , 2	$2-2=0$ , 1 , 1	$2-1=1$ , 1 , 1		

Total Allocations = 6.

$$T.C = 1 \times 10 + 1 \times 20 + 2 \times 10 + 3 \times 10 + 2 \times 30 + 2 \times 20 = \text{Rs } 180.$$

Unsolved Questions on VAM:

Q1)

	$d_1$	$d_2$	$d_3$	$d_4$	Supply
$O_1$	2	3	11	7	6
$O_2$	1	0	6	1	1
$O_3$	5	8	15	9	10
Reqd.	7	5	3	2	

Q2)

		Warehouses				
		$W_1$	$W_2$	$W_3$	$W_4$	Capacity
factory	$F_1$	19	30	50	10	7
	$F_2$	70	30	40	60	9
	$F_3$	40	8	70	20	18
	Reqd.	5	8	7	14	



6) MODI METHOD  $\rightarrow$  to find optimal solution. (6)

Q)

	$W_1$	$W_2$	$W_3$	Supply
$F_1$	16	20	12	200
$F_2$	14	8	18	160
$F_3$	26	24	16	90
Demand	180	120	150	

Step 1) Apply VAM to find most feasible soln.

	$W_1$	$W_2$	$W_3$	Supply	Row Penalties
$F_1$	<del>16</del> <sup>140</sup>	20	<del>12</del> <sup>60</sup>	<del>200</del> <sup>140</sup>	4, 4, 4
$F_2$	<del>14</del> <sup>40</sup>	<del>8</del> <sup>120</sup>	18	160	6, 4, 4
$F_3$	26	24	<del>16</del> <sup>90</sup>	90	8, 10 $\leftarrow$
Demand	<del>180</del> <sup>40</sup>	<del>120</del> <sup>120</sup>	<del>150</del> <sup>60</sup>		
Column Penalties	2, 2, 2, 2	$\uparrow$ 2	4, 4, 6		

Check whether it is a basic feasible soln. or not.

Reqd. no. of allocations =  $m+n-1 = 5$ .  
 & actual no. of allocations = 5.

Step 2:) Assign dual variables:

Check  $\leq$  how or column has maximum no. of allocations. Assign that row or column, a dual variable = 0.

Let  $U_i$  = no. of rows  
 $V_j$  = no. of columns.

Using formula,  $U_i + V_j = C_{ij}$

	$W_1$	$W_2$	$W_3$	$\Sigma$	Dual variables $U_i$
$F_1$	16 <sup>140</sup>	20 <sup>-10</sup>	12 <sup>60</sup>	200	$U_1 = 0$
$F_2$	14 <sup>40</sup>	8 <sup>120</sup>	18 <sup>-8</sup>	160	$U_2 = -2$
$F_3$	26 <sup>-6</sup>	24 <sup>-10</sup>	16 <sup>90</sup>	90	$U_3 = 4$
Reqd.	180	120	150		
O.V, $V_j$	$V_1 = 16$	$V_2 = 10$	$V_3 = 12$		

Assign  $U_1 = 0$  as Row 1 is having 2 allocations.  
 $\therefore U_1 + V_1 = 16$  ; Similarly find other dual variables.  
 $\Rightarrow V_1 = 16$

Step 3) Calculate Opportunity cost to other non-basic cells (non-allocated cells), using the formula,  $U_i + V_j - C_{ij} = O.C$



(5)

$$u_1 + v_2 - 20 = -10 \quad \left\{ \begin{matrix} 20 \\ -10 \end{matrix} \right\}$$

$$u_2 + v_3 - 18 = -8$$

$$u_3 + v_1 - 26 = -6$$

$$u_3 + v_2 - 24 = -10$$

(7)

Since all the values of O.C. are negative, we need not to reduce further cost. We have reached the optimal solution.

$$\therefore \text{Total Cost} = 16 \times 140 + 12 \times 60 + 14 \times 40 + 8 \times 120 + 16 \times 90 \\ = \text{Rs } 5920/-$$

Unsolved:

	$d_1$	$d_2$	$d_3$	$d_4$	Supply
$S_1$	23	27	16	18	30
$S_2$	12	17	20	51	40
$S_3$	22	28	12	32	53
Reqd.	22	35	25	41	

Cases of MODI METHOD:

- Degeneracy Case
- Unbalanced Matrix  $\rightarrow$  Dummy Line
- 0-allocations.